

Supplementary material: Proofs of Lemmas 1 and 3

Proof of Lemma 1. If $M = m$ but $\widehat{M} \neq m$, then

$$\|\widehat{f}_n - f_m\|_{L_1(\Pi)} \geq \|f_m - \widehat{f}_M\|_{L_1(\Pi)} - \|\widehat{f}_n - \widehat{f}_M\|_{L_1(\Pi)} \geq 2\varepsilon - \|\widehat{f}_n - \widehat{f}_M\|_{L_1(\Pi)}. \quad (28)$$

Thus, if $\|\widehat{f}_n - f_m\|_{L_1(\Pi)} < \varepsilon$, then it must be the case that $\|\widehat{f}_n - \widehat{f}_M\|_{L_1(\Pi)} < \varepsilon$, which, in view of (28), is a contradiction. Hence,

$$\mathbb{P}^{S,\pi}(\widehat{M} \neq M) \leq \mathbb{P}^{S,\pi}(\widehat{f}_n \notin \mathcal{F}_\varepsilon(f_M)) \leq \delta,$$

and the lemma is proved. \square

Proof of Lemma 3. We only prove (16), since the proof of (17) is similar. First note that

$$\frac{\mathbb{P}^{S,m}(x^n, y^n)}{\mathbb{Q}^S(x^n, y^n)} = \prod_{t=1}^n \frac{P_{Y|X}^m(y_t|x_t)}{Q_{Y|X}(y_t|x_t)} = \prod_{x \in \mathcal{X}} \prod_{t: x_t=x} \frac{P_{Y|X}^m(y_t|x)}{Q_{Y|X}(y_t|x)}.$$

Then

$$\begin{aligned} D(\mathbb{P} \parallel \mathbb{Q}) &= \frac{1}{N} \sum_{m=1}^N \sum_{x^n, y^n} \mathbb{P}^{S,m}(x^n, y^n) \log \frac{\mathbb{P}^{S,m}(x^n, y^n)}{\mathbb{Q}^S(x^n, y^n)} \\ &= \frac{1}{N} \sum_{m=1}^N \sum_{x \in \mathcal{X}} \mathbb{E}_{\mathbb{P}^{S,m}} \left[\sum_{t: X_t=x} \log \frac{P_{Y|X}^m(Y_t|x)}{Q_{Y|X}(Y_t|x)} \right] \\ &= \frac{1}{N} \sum_{m=1}^N \sum_{x \in \mathcal{X}} D(P_{Y|X}^m(\cdot|x) \parallel Q_{Y|X}(\cdot|x)) \mathbb{E}_{\mathbb{P}^{S,m}} [N(x|X^n)], \end{aligned}$$

which gives (16). Eq. (18) follows from the fact that, for a passive strategy, the expectation of $N(x|X^n)$ is equal to $n\Pi(x)$ under both $\mathbb{P}^{S,m}$ and \mathbb{Q}^S . The same proof holds with D_{re} replaced by D . \square

Proof of Lemma 4. The proof is via the probabilistic method. Specifically, we shall show that if we select N binary strings from $\{0, 1\}_d^k$ uniformly at random, then the resulting set will have all three desired properties with probability strictly greater than 0.

For a fixed $\beta \in \{0, 1\}_d^k$, let $\mathcal{U}_\beta \triangleq \{\beta' \in \{0, 1\}_d^k : d_H(\beta, \beta') \leq d\}$. Then for any $\beta' \in \mathcal{U}_\beta$ $|\{i \in [k] : \beta_i = \beta'_i = 1\}| \geq \frac{d}{2}$. Hence,

$$|\mathcal{U}_\beta| \leq \binom{d}{d/2} \binom{k-d/2}{d/2} \leq \binom{d}{d/2} \binom{k}{d/2} \leq 2^d \binom{k}{d/2},$$

where we have used the fact that $\binom{d}{\ell} \leq 2^d$ for any $\ell \leq d$. From this we see that if we draw an element of $\{0, 1\}_d^k$ uniformly at random, then it will be in $|\mathcal{U}_\beta|$ with probability

$$p = \frac{|\mathcal{U}_\beta|}{|\{0, 1\}_d^k|} \leq \frac{2^d \binom{k}{d/2}}{\binom{k}{d}}.$$

Thus, if we select N elements of $\{0, 1\}_d^k$ uniformly at random, then the probability that the j th element will be d -close in the Hamming distance to any of the $j-1$ already selected ones is at most $(j-1)p$, and the probability that any two elements are d -close is at most $(N^2/2)p$. Hence, with the choice $N = \lfloor (3k/16d)^{d/4} \rfloor \geq (k/6d)^{d/4}$ we have

$$\frac{N^2 p}{2} \leq \frac{1}{2} \left(\frac{3k}{16d} \right)^{d/2} \frac{2^d \binom{k}{d/2}}{\binom{k}{d}} \leq \frac{1}{2},$$

where we have used the fact that $\binom{k}{d}/\binom{k}{d/2} \geq \left(\frac{k}{d} - \frac{1}{2}\right)^{k/2}$, as well as the fact that $\frac{3k}{4d} \leq \frac{k}{d} - \frac{1}{2}$ for $d \leq k/2$. Hence, with probability at least $1/2$, all the N elements will be strictly d -separated.

We now show that the randomly selected set of N elements of $\{0, 1\}_d^k$ will also be “well-balanced” in the sense of (19) with probability strictly larger than $1/2$. To that end, let us fix $j \in [k]$ and let Z_1, \dots, Z_N be the $\{0, 1\}$ -valued random variables, corresponding to the j th coordinates of the randomly chosen elements. Observe that $\mathbb{E}Z_i = d/k$. Then Bernstein’s inequality gives

$$\begin{aligned} \Pr \left(\left| \frac{1}{N} \sum_{i=1}^N Z_i - \frac{d}{k} \right| > \frac{d}{2k} \right) &\leq 2 \exp \left(- \frac{N(d/2k)^2}{2(d/k)(1 - d/k) + 2(1 - d/k)(d/(2k))/3} \right) \\ &= 2 \exp \left(- \frac{Nd}{12k} \right) \end{aligned}$$

This, together with the union bound, shows that the probability of (19) being violated is at most $2k \exp \left(- \frac{Nd}{12k} \right)$, which will be strictly less than $1/2$ for sufficiently large d . Hence, the probability that a set of N elements of $\{0, 1\}_d^k$ drawn uniformly at random will fail to satisfy either the separation condition (ii) or the balance condition (iii) is strictly less than 1. This completes the proof. \square